## Book

## A Simplified Approach to

## Data Structures

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## B TREE

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## B-tree

A very important category of m-way trees is b-tree which was introduced by R.Bayer and E.Mcgreight B-tree of order $m$ has following properties:

- Each node of tree except root node must have maximum of $\mathbf{m - 1}$ keys and minimum of $\mathbf{m} / \mathbf{2 - 1}$ keys.
- All leaf nodes in B-tree must be on same level.
- Each node of tree except root node and leaf nodes has maximum of $\mathbf{m}$ children and minimum of $\mathbf{m} / \mathbf{2}$ children.
- The key values in each node are stored in ascending order.
- The key in node separates the ranges of keys which are atored in each child of that node. A B-tree of order $m$ and height $h$ has total number of elements $=\mathbf{m}^{\wedge}(h+1)-1$.


## Operations performed on B-tree.

1. Searching
2. Insertion
3. Deletion

- Searching:
if the B-tree is of order 4 then we have maximum of 4 children of each node and hence 4 way choice to move to one of 4-children.the correct child is chosen by performing a linear search of keys in a node.


## SEARCHING IN B TREE

Value to find is 145
Which may be between 100-150 in parent node X

## For example, consider a B tree of order 4.

Suppose we want to search 145 then first of all we go to root node and find that it is between 100-150, so we go to third child of this root node. In this node we again perform linear search.
This time we find the desired key value in this node

## Example of Insertion in B-Tree(1)



## Example of Insertion in B-Tree(2)



Insert <60, 80, 35, 75, 43, 32, 78, 65, 55, 68, 99, 92, 96, 98, 94>

Example of Insertion in B-Tree(3)


## Delete 2



Now, we delete 2

## Delete 21



## Delete 10



Deleting 10 causes node c to underflow. This causes the parent, node g to recombine with nodes f and $a$. This causes the tree to shrink one level.

## Delete 10



Deleting 10 causes node c to underflow. This causes the parent, node $g$ to recombine with nodes $f$ and a. This causes the tree to shrink one level.

